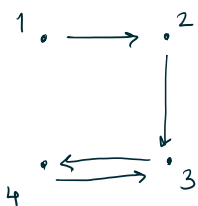


\* Last time : Adjacency matrices

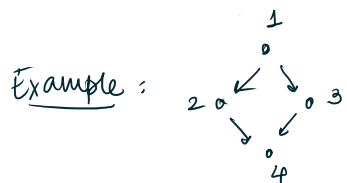
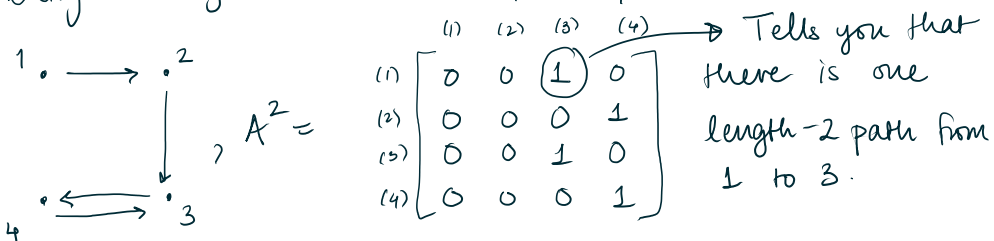
\* Today's topic : Computing with adjacency matrices!



$$A = \begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

$$A^2 = \begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix} \begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix} = \begin{bmatrix} 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

Tip: Ex. In 2<sup>nd</sup> row, the third entry is 1 & everything else is 0. This third entry will multiply with the elements in the 3<sup>rd</sup> row of the second copy of A, so in this case we only need to look at the 3<sup>rd</sup> row of A, multiplied with this entry  $\rightarrow$  to get the 2<sup>nd</sup> row of the product.



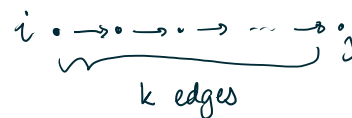
$$A = \begin{bmatrix} 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$A^2 = \begin{bmatrix} 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix} \begin{bmatrix} 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

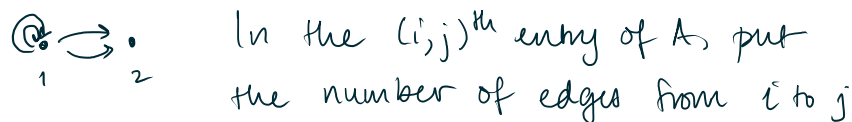
$$= \begin{bmatrix} 0 & 0 & 0 & 2 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$0 \cdot 0 + 1 \cdot 1 + 1 \cdot 1 + 0 \cdot 0$   
 $\downarrow \rightarrow 1$     $\underbrace{1 \rightarrow 2}$     $\underbrace{1 \rightarrow 3}$     $\underbrace{1 \rightarrow 4}$   
 $\downarrow \rightarrow 4$     $\underbrace{2 \rightarrow 4}$     $\underbrace{3 \rightarrow 4}$     $\underbrace{4 \rightarrow 4}$

Thm: The  $(i, j)$ <sup>th</sup> entry of  $A^k$  is exactly the number of length- $k$  paths from  $i$  to  $j$



\* Aside: If you had a non-simple graph:



$$A = \begin{bmatrix} 2 & 2 \\ 0 & 0 \end{bmatrix}$$

$$A^2 = \begin{bmatrix} 2 & 2 \\ 0 & 0 \end{bmatrix} \begin{bmatrix} 2 & 2 \\ 0 & 0 \end{bmatrix} = \begin{bmatrix} 4 & 4 \\ 0 & 0 \end{bmatrix}$$

(Variant of the adjacency matrix, but the theorem still works!)

Rmk: We can now answer the question of finding length- $k$  paths between vertices.

\* Connectedness?

Q: Is there always at least one path from  $i$  to  $j$ , of some length? (For some given vertices  $i$  &  $j$ ?)

Ans You can look at the  $(i,j)^{th}$  entry of

$(A + A^2 + A^3 + \dots)$  add corresponding entries,

& if this entry ever becomes non-zero, then you know the answer is yes.

If the graph has  $n$  vertices, and if there is some path from  $i$  to  $j$ , you're guaranteed at least one path of length at most  $(n-1)$ .

$(A + A^2 + A^3 + \dots + A^{n-1})$  In this sum, you can stop after  $A^{n-1}$ .

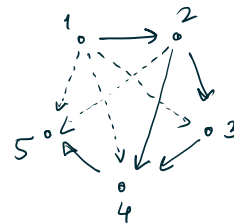
If you don't see a path from  $i$  to  $j$  until  $A^{n-1}$ , you will never see one.

This is because the longest <sup>possible</sup> length of the shortest path from  $i$  to  $j$  is  $(n-1)$



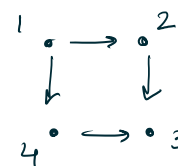
Exercise: Do some examples!

Transitive closure of a relation?



Transitive closure obtained by adding the dotted edges.

In terms of matrices:  $\rightarrow (1,3)$  needs to be an edge.



$$A = \begin{bmatrix} 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

(the graph of a relation)

$$A^2 = \begin{bmatrix} 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

there is a length 2 path from 1 to 3!

For transitive closure: If you get a positive entry in  $A^k$  at spot  $(i,j)$ , then add an edge from  $i$  to  $j$ .

Algebraically: We encode this procedure by

"boolean matrix product."

Next Wednesday