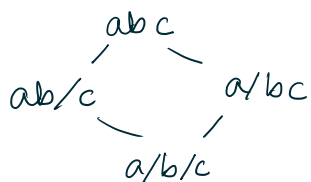


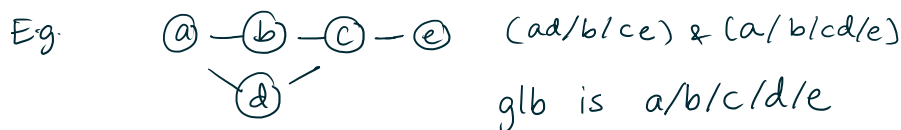
* Recap: Given a graph (V, E) [undirected w/o self loops] we can write down its bond poset:

A bond $B = \{V_1, \dots, V_k\}$ is \leq a bond $C = \{W_1, \dots, W_\ell\}$ if every V_i is entirely contained within a single W_j .
Equivalently, each W_j is a union of some number of pieces V_i .

* Example: $(a) - (b) - (c)$

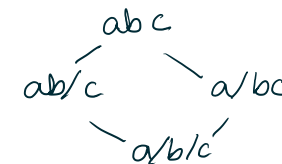


* Rmk: In fact, the bond poset is a lattice.
Exercise: Figure out the lub/glb in this poset

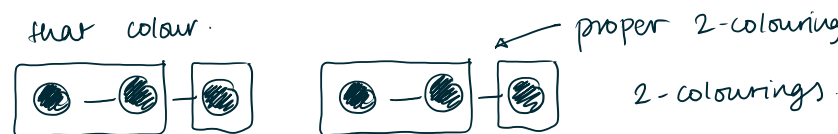


lub? ad & cd must be clubbed & ce is also clubbed
 " $acde/b$

E.g. $(a) - (b) - (c)$



Recall: A k -colouring of a bond $B = \{V_1, \dots, V_k\}$ means that you assign one of k different colours to each V_i , such that all vertices in V_i are coloured with that colour.



A proper k -colouring is one where if V_i & V_j are connected by an edge, they get different colours.

* Notation: Let B_0 be the smallest bond
 $\{ \{v\} \mid v \in V \}$

Note: A proper k -colouring of B_0 is exactly a proper k -colouring of the graph G .

Defn: Let $P_t(B) =$ number of t -colourings of a bond B
 $q_t(B) =$ number of proper t -colourings of a bond B .

We want to compute $q_t(B_0) = \#$ of proper t -colourings of the graph G .

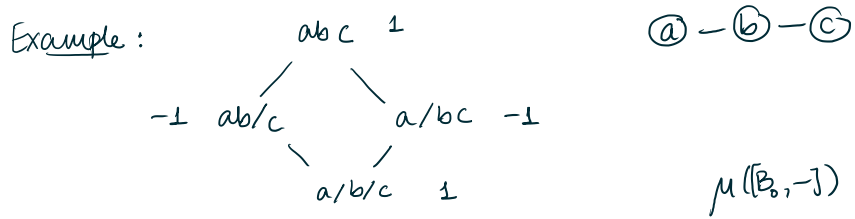
Note: $P_t(B) = t^{\# \text{parts in } B}$

And: $P_t(B) = \sum_{B \geq C} q_t(C)$ convince yourself!

$$P_t(B) = (\sum * q_t)(B)$$

Möbius inversion $\Rightarrow q_t(B) = (\mu * P_t)(B)$

$$q_t(B_0) = \sum_C \mu([B_0, C]) \cdot P_t(C)$$



$$q_t(B_0) = 1 \cdot t^3 - 1 \cdot t^2 - 1 \cdot t^2 + 1 \cdot t = t^3 - 2t^2 + t = t(t-1)^2$$

Theorem: The chromatic polynomial of G has the following formula:

$$P_G(t) = \sum_{C \text{ bond}} \mu([B_0, C]) \cdot P_t(B)$$

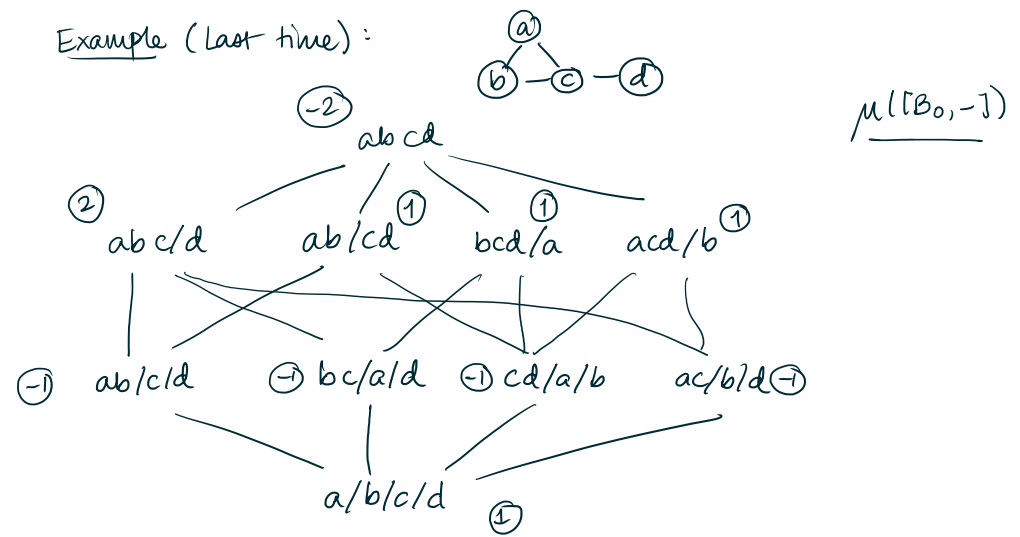
$$= \sum_{C \text{ bond}} \mu([B_0, C]) \cdot t^{(\# \text{ parts of } B)}$$

* Remark: This method also proves that $P_G(t)$ is a polynomial!

Recall: $\mu([x, x]) = 1 \quad \forall x$

$$\mu([x, y]) = -\sum_{x \leq z \leq y} \mu([x, z])$$

Example (last time):



$$P_G(t) = t^4 - 4t^3 + 5t^2 - 2t = t(t-1)^2(t-2)$$

you can check directly in this case

Exercise: Do this for the graph in the worksheet!

* Done with posets for this course !!

"Machines" — Regular expressions & finite automata.

A way of pattern-matching, for example for English words.

Defn: An alphabet is a finite set of symbols; typically denoted Σ . (sigma)

E.g. $\Sigma^1 = \{a, b, c\}$ or $\Sigma^1 = \{0, 1\}$

$$\Sigma = \{a, b, c, \dots, z\}$$

Defn: A word in an alphabet Σ is a finite ordered list of elements in Σ , written without punctuation next to each other.

Eg. if $\Sigma = \{0, 1\}$ then 000, 111, 101011 are words

The "empty word" is also a word; it is denoted ϵ .
(So as a convention, Σ should not contain the symbol ϵ)

The set of all possible words of Σ , including the empty word, is denoted Σ^* .
(Σ^* is infinite unless $\Sigma = \emptyset$; in that case, $\Sigma^* = \{\epsilon\}$)

Defn: A language L is just a subset of Σ^* .
(L tells you which words are valid words in that language)