MATH 2301 30 Sept 2020

* Recap: We defined valid regular expressions over an alphabet Σ.

Precedence order: $* \gg$ concatenation $\gg |$

(In the absence of parentheses, apply $*$ first, concat next and $|$ last.

Just like in an algebraic expression:

$\wedge \gg \times \gg +$)

* We have $r \longleftrightarrow L(r) \subseteq \Sigma^*$

the language of r

* We say $s \in \Sigma^*$ matches r if $s \in L(r)$.

* Theoretical vs practical regular expressions: a demo

[We saw $+$, $\cdot$, complement $\sim$ extra syntax in practical regex]

We exclude this syntax and keep minimal syntax so it's easier to prove theorems. Many of the extra operations can be simulated with the ones we have]

* Today: Which languages can be expressed as $L(r)$ for some regex r?

Claim: The language $L = \{0^n 1^n \mid n \in \mathbb{N}\}$

(n 0s followed by n 1s)

is not equal to $L(r)$ for any regex r.

Exercise: Try (and fail) to construct a regex for L.

We'll try to prove this claim later.

Def: A language L is called regular if $L = L(r)$ for some regex r.

[Not all languages will be regular.]

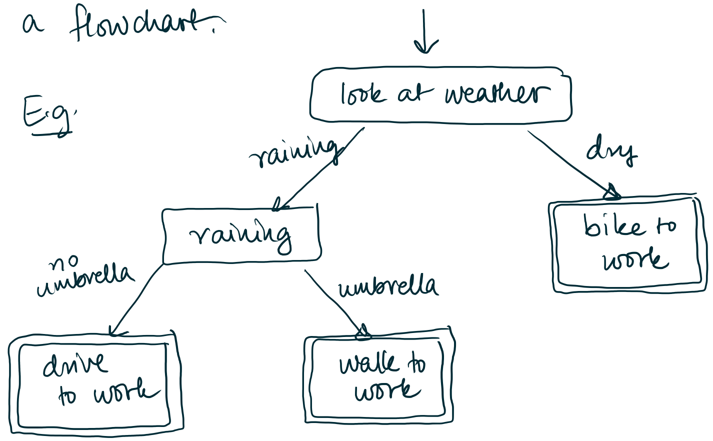
* (Deterministic) finite automata, or DFAs

(Finite state machines)

These will end up being machines/model mini computers that exactly recognise regular languages.

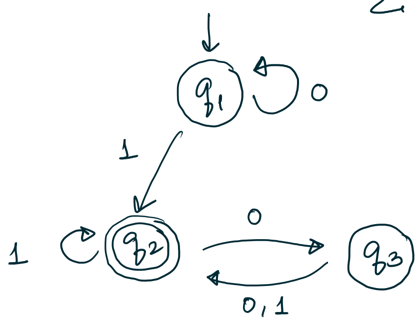
Informally, a DFA is a machine that simulates a flowchart.

E.g.



E.g.

Here is a DFA: [you need an alphabet Σ, say $\Sigma = \{0,1\}$]

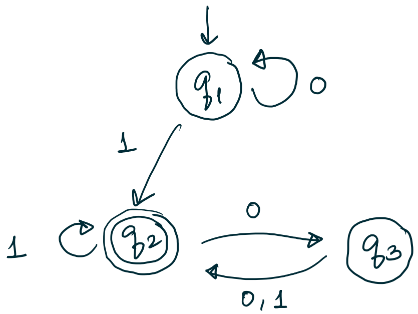


States: q_1, q_2, q_3

Start state: q_1

Each state has one outgoing arrow labelled by each $a \in \Sigma$.

Every arrow except for the start arrow is labelled by some $a \in \Sigma$.

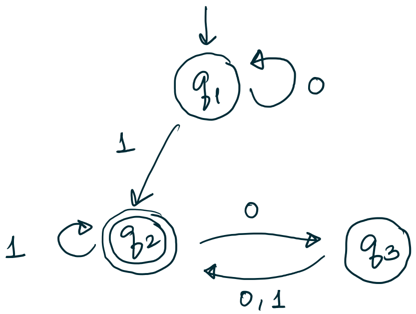


(double circled)
q_2 is an <u>accept state</u>

q_1, q_3 are <u>reject states</u>
(or fail / failure states)

Given any word $w \in \Sigma^*$ you can "run" the DFA on that word, and the DFA will either accept or reject w.

<u>Eg.</u> $w = 011101$
Start at q_1. "Read" the first letter, follow the corresponding arrow, continue.



Start at q_1
1. Read 0 $\to$ go to q_1
2. Read 1 $\to$ go to q_2
3. Read 1 $\to$ go to q_2
4. Read 1 $\to$ go to q_2
5. Read 0 $\to$ go to q_3
6. Read 1 $\to$ go to q_2

Since q_2 is an accept state, we <u>accept</u> w.

$w = 10, 010, 0110, 0, \varepsilon$ are all rejected.

Suggestion (Mikayla) : $0^* 1 (1 | 00 | 01)^* = r$
then $L(r)$ is exactly the set of accepted strings of this DFA (Check!)