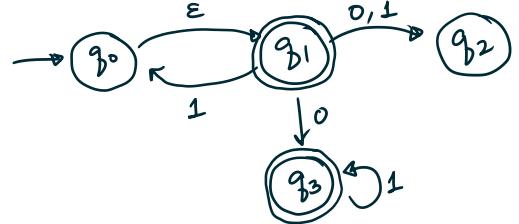


MATH 2301

* Non-deterministic finite automata (NFA).

** Informal description



It is almost exactly like a DFA, with the following key differences:

- i) Arrows may be labelled by ϵ
- ii) Each state can have zero or more outgoing arrows labelled by each letter of Σ .

** Note : An NFA has only one start state, and any number of accept states.

** How to compute with an NFA?

Given any string w , try all possible ways to run w through the NFA.

If at least one option reaches an accept state, and you have read all of w , then ACCEPT w . Otherwise, REJECT w .

** Formal definition

Fix an alphabet Σ

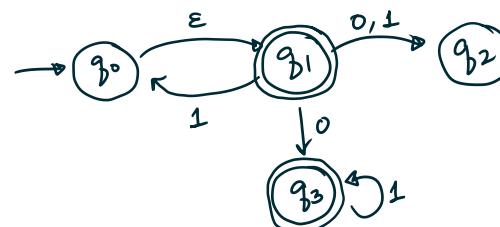
*** Def : A nondeterministic finite automaton (NFA) consists of the following:

- 1) A set of states Q .
- 2) A single start state $q_0 \in Q$.
- 3) A set of accept states $A \subseteq Q$.
- 4) A transition function

$$\Delta : Q \times (\Sigma \cup \{\epsilon\}) \rightarrow P(Q)$$

↑ ↑
 state the letter
 you're at (or ϵ) that
 you read

↑ power set of Q .



Example values of Δ

$$\Delta(q_0, 1) = \emptyset \in P(Q)$$

$$\Delta(q_0, 0) = \emptyset$$

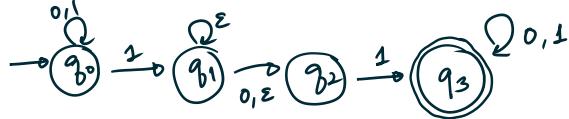
$$\Delta(q_0, \epsilon) = \{q_1\}$$

$$\Delta(q_1, 0) = \{q_2, q_3\}$$

$$\Delta(q_3, 1) = \{q_3\}$$

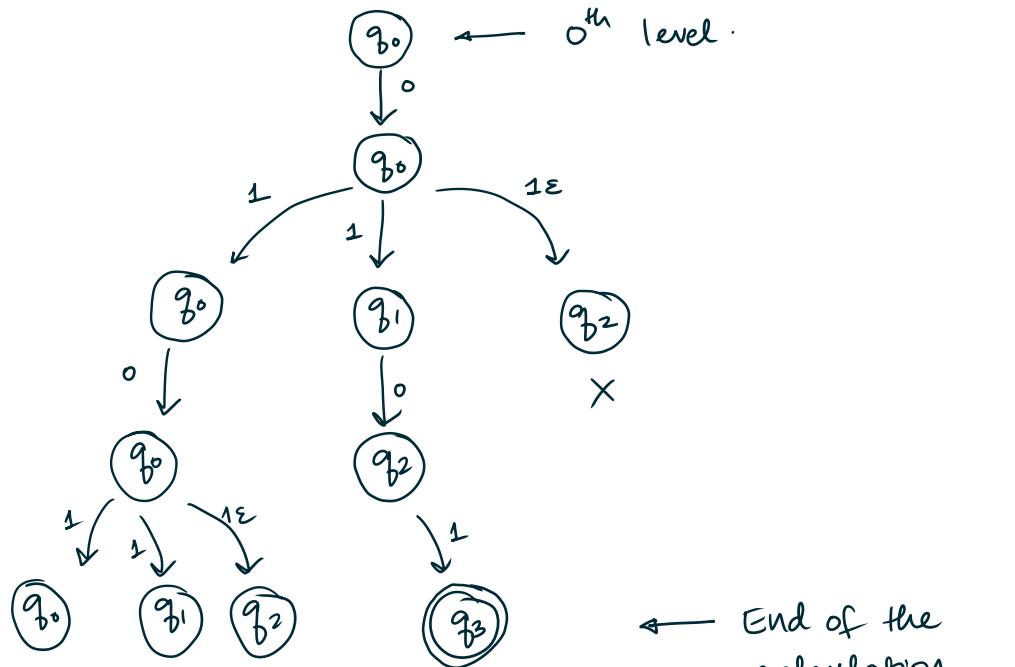
$$\Delta(q_2, \epsilon) = \emptyset$$

**** Example**



$w = 0101$

Draw a calculation tree :



Result : ACCEPT, because q_3 is reachable at level 4, and is accepting.

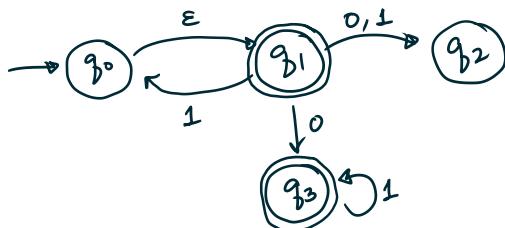
**** Key points about the calculation / calculation tree**

- When you read a letter, you should also read any ϵ s that come before & after.

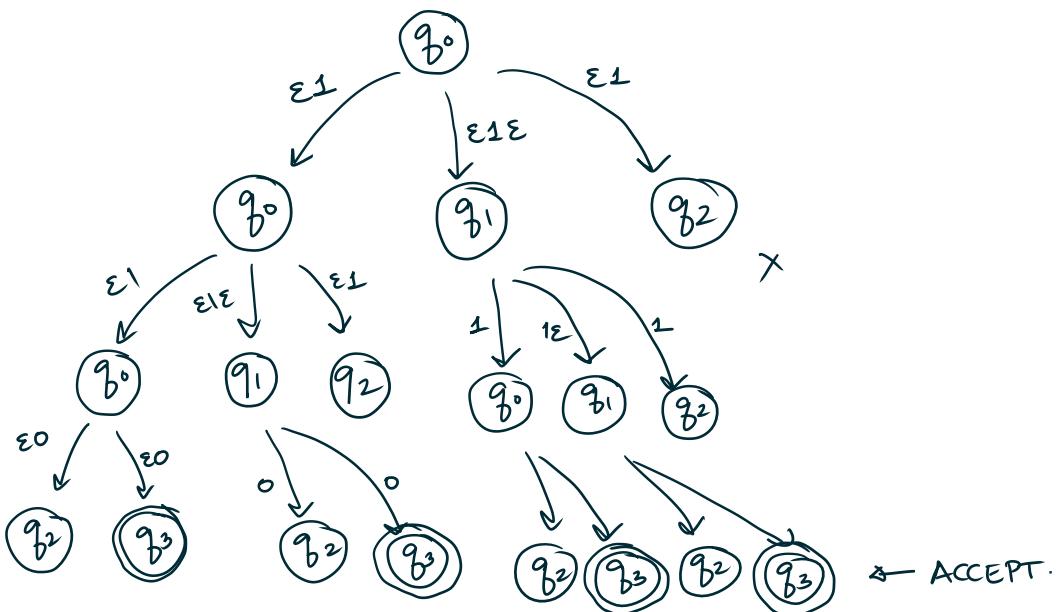
E.g. If you're reading "1", you should

follow all paths of the form $\underbrace{\epsilon \epsilon \dots \epsilon}_\text{any number of } \epsilon 1 \underbrace{\epsilon \dots \epsilon}_\text{any number of } \epsilon$

- self-loops labelled ϵ can be ignored
- Branches that don't reach the last level (i.e. the end of the string) are failed branches.
- If none of the branches reach the last level, reject.



$w = 110$



→ ACCEPT.