

## MATH 2301

### \* Nim ordering classification

You can consider the heaps to be un-ordered.

But keep multiplicity!

E.g.  $(1, 1, 2) = (2, 1, 1) = (1, 2, 1)$  but  
 $(1, 1, 2) \neq (1, 2)$ .

### \* Nim strategy

#### \*\* Experiments with XOR

XOR = exclusive OR, is an operation on binary strings, defined as follows.

$$\begin{array}{r} \text{Examples} \quad 101_2 \\ \oplus \quad 101_2 \\ \hline 110_2 \end{array}$$

Def: The XOR operation is defined as

$$\{0, 1\} \times \{0, 1\} \rightarrow \{0, 1\}$$

$$0 \oplus 1 := 1$$

$$1 \oplus 0 := 1$$

$$0 \oplus 0 := 0$$

$$1 \oplus 1 := 0$$

(aka addition mod 2)

Extend this to binary strings by right-aligning them and applying XOR column-wise.

(If one string is shorter than the other, then pad it with 0s on the left.)

$$\begin{array}{r} \text{E.g.} \quad 111_2 \\ \quad \quad 01_2 \\ \oplus \quad 11011_2 \\ \hline 11101_2 \end{array}$$

Rmk

- XOR is commutative:

$$x \oplus y = y \oplus x$$

- XOR is associative:

$$(x \oplus y) \oplus z = x \oplus (y \oplus z)$$

Rmk:

Two binary strings are considered equal if they are the same up to an initial sequence of zeroes:

$$01_2 = 1_2 = 001_2$$

-  $x \oplus x = 0$  for every  $x$ .

#### \*\* Some nim games (Nim sum)

-  $(3, 3, 5)$ : convert each pile size to a binary number: write as the sum of distinct powers of 2, and put a 1 for every power that appears, and a 0 for every power that doesn't appear.

$$3 = 2 + 1 = 2^1 + 2^0 = 11_2$$

$$1 \quad 1$$

$$5 = 4 + 1 = 2^2 + 2^0 = 101_2$$

Nim sum of (3, 3, 5) is also denoted

$$3 \oplus 3 \oplus 5 = \underbrace{11_2 \oplus 11_2}_{=0} \oplus 101_2 = 101_2$$

---

(2, 3, 1)

$$2 \oplus 3 \oplus 1 = 10_2 \oplus 11_2 \oplus 1_2 = \begin{array}{r} 10 \\ 11 \\ \oplus 01 \\ \hline 00_2 \end{array} = 0_2$$

$$(n, n) \text{ and } n \oplus n = 0_2$$

\*\* Def: Let  $(n_1, \dots, n_k)$  be a game state for nim. Its nim-sum is  $n_1 \oplus \dots \oplus n_k$ . That is, the XOR of the binary representations of  $n_1, n_2, \dots, n_k$ .

\*\* Theorem: A game position in nim is a:  
 P position iff its nim-sum is 0.  
 N position iff its nim-sum is non-zero.

\*\* Example: (4, 5, 6, 13)

$$\begin{array}{r} 100_2 \\ 101_2 \\ 110_2 \\ \hline 1101_2 \\ 1010_2 \end{array} \left. \vphantom{\begin{array}{r} 100_2 \\ 101_2 \\ 110_2 \\ \hline 1101_2 \\ 1010_2 \end{array}} \right\} \Rightarrow \text{winning position (by thm)} \quad (N)$$

Theorem + our knowledge of N/P labeling tells us that there must be a move that takes us to a P position, i.e. one whose nim-sum is zero.

\*\* Example: (4, 5, 6, 13)  $\xrightarrow{?}$  (4, 5, 6, 7)

$$\begin{array}{r} 100_2 \\ 101_2 \\ 110_2 \\ \hline 1101_2 \\ 1010_2 \end{array} \qquad \begin{array}{r} 100_2 \\ 101_2 \\ 110_2 \\ \hline 111_2 \\ \hline 000_2 \end{array}$$

A winning move is one that keeps an even number of 1s in each column  $\Leftrightarrow$  XOR is zero

\*\* Example (7, 4, 5)  $\rightarrow$  (1, 4, 5)

$$\begin{array}{r} 111_2 \\ 100_2 \\ \oplus 101_2 \\ \hline 110_2 \end{array} \begin{array}{l} \rightarrow (7, 2, 5) \\ \rightarrow (7, 4, 3) \end{array}$$

Observe: A state is "N" iff <sup>at least</sup> one of the columns has an odd number of 1s.

Strategy: Consider the leftmost column w/ an odd number of 1s. Choose any one number that has a 1 in this column.

$$\begin{array}{r} \textcircled{111_2}^n \\ 100_2 \\ \oplus 101_2 \\ \hline \textcircled{110_2}^s \end{array}$$

Strategy: Take this number,  
say  $n$ .

Let  $s$  be the nim-sum

Take  $n \oplus s$ . Convert to decimal,  
and replace  $n$  by this new  
number.

$$\begin{array}{r} 111_2 \\ \oplus 110_2 \\ \hline \end{array}$$

$$001_2 = 1 \text{ (in decimal)} \Rightarrow (7, 4, 5) \rightarrow (1, 4, 5)$$

$$\text{or, } \begin{array}{r} 100 \\ \oplus 110_2 \\ \hline 010_2 = 2 \end{array}$$

$$\Rightarrow (7, 4, 5) \rightarrow (7, 2, 5)$$

$$\text{or } \begin{array}{r} 101_2 \\ \oplus 110_2 \\ \hline 011_2 = 3 \end{array}$$

$$\Rightarrow (7, 4, 5) \rightarrow (7, 4, 3)$$