

* Today : Nim

Game state = some number of piles/heaps, consisting each of some number (≥ 0) of berries.

(Order is irrelevant, but keep track of multiplicity)

A move consists of eating a non-zero number of berries from exactly one pile.

E.g. : $(1, 2, 2) \xrightarrow{\text{any move}} (0, 2, 2), (1, 1, 2), (1, 0, 2), \dots$ (game states are tuples) ~~and ordering~~

The game ends when there are no moves.

* Easy cases

- If \exists exactly one pile of $\neq 0$ berries, then the first player wins.
- \Rightarrow Any position of the form (m) , for ~~where~~ $m > 0$, is an N-position.

- Positions of the form (a, b) for $a, b > 0$

$$\textcircled{N} (5, 8) \quad \text{or} \quad (4, 4) \textcircled{P}$$

$$(4, 4) \xrightarrow{\downarrow} (2, 4) \xrightarrow{\downarrow} (2, 2) \xrightarrow{\dots}$$

* Claim : A position of the form (a, b) where $a, b \geq 0$ and $a \neq b$ is an N-position. A position of the form (a, b) with $a = b$ is a P-position.

Pf (sketch) : If (a, b) is such that ~~where~~ $a \neq b$ and $a, b \geq 0$, then, suppose (WLOG) that $a > b$. The winning move is

$$(a, b) \xrightarrow{\text{move}} (b, b)$$

Subsequently, no matter what the second player does, the first player can "mirror" the move.

$$\text{E.g. } (4, 7) \rightarrow (4, 4) \rightarrow (1, 4) \rightarrow (1, 1)$$

Similarly, from a position of the form (a, a) , the only possible moves go to positions of the form (c, d) with $c \neq d$.

* Lesson we learn : "Mirroring" is a good strategy to keep in mind.

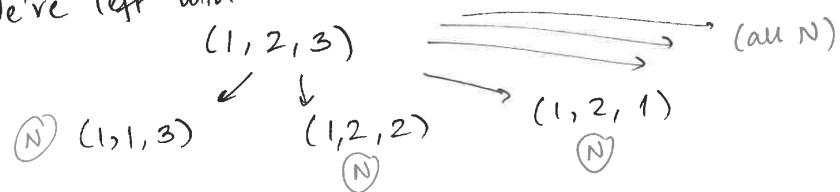
Q: What about starting states with > 2 non-empty piles?

(3)

E.g. (1, 2, 3)

observe: It's not in P1's favour to eat any single pile.

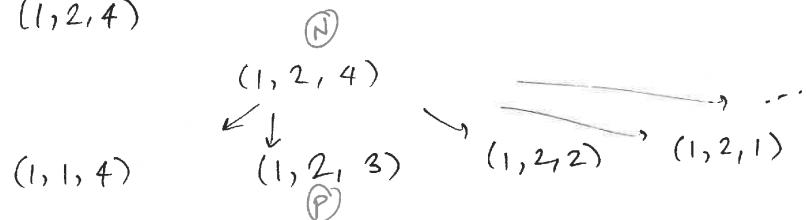
We've left with:



→ (1, 2, 3) is \textcircled{P} .

(Note: (1, 1, 3) is \textcircled{N} because there is a move to (1, 1, 0), which is \textcircled{P})

E.g. (1, 2, 4)



(1, 2, 4) is \textcircled{N} because it has a move to (1, 2, 3), which is \textcircled{P} .

Let's work towards a solution

Given a position (m_1, m_2, \dots, m_k) we're going to compute its "nim-sum": $m_1 \oplus m_2 \oplus \dots \oplus m_k$.

= decimal corresponding to the XOR of the binary representation of $m_1, m_2, m_3, \dots, m_k$.

The nim-sum is defined in terms of XOR
= exclusive or.

Given some m , we can convert it to binary.

$$\text{E.g. } m = 5 = 1 \cdot 2^2 + 0 \cdot 2^1 + 1 \cdot 2^0 = 2^2 + 1$$

$$5 = (101)_2$$

Expand m as a sum of distinct powers of 2,
something like $a_k \cdot 2^k + a_{k-1} \cdot 2^{k-1} + a_{k-2} \cdot 2^{k-2} + \dots + a_0 \cdot 2^0$
where each a_i is either 1 or 0.

$$\text{Then } m = (a_k a_{k-1} \dots a_0)_2$$

$$\text{E.g. } 19 = 16 + 2 + 1 = 1 \cdot 2^4 + 0 \cdot 2^3 + 0 \cdot 2^2 + 1 \cdot 2^1 + 1 \cdot 2^0$$

$$19 = (10011)_2$$

E.g. (nim sum) of (5, 7, 19)

5 =	101_2	(column-wise XOR)
\oplus 7 =	111_2	$\begin{bmatrix} 1 \oplus 1 = 0 \\ 1 \oplus 0 = 1 \\ 0 \oplus 1 = 1 \\ 0 \oplus 0 = 0 \end{bmatrix}$
\oplus 19 =	$\begin{array}{r} 10011_2 \\ \hline 10001_2 \end{array}$	