

** Last time : Nim

** Today : Winning strategy (?) using nim-sum.

Recall: The nim-sum of (m_1, \dots, m_k) is the (decimal representation of) the column-wise XOR of the binary representations of m_1, m_2, \dots, m_k .

Example : $(4, 9, 11)$

$$4 = 2^2 = 1 \cdot 2^2 + 0 \cdot 2^1 + 0 \cdot 2^0 = 100_2$$

$$9 = 8 + 1 = 2^3 + 1 = 1 \cdot 2^3 + 0 \cdot 2^2 + 0 \cdot 2^1 + 1 \cdot 2^0 = 1001_2$$

$$11 = 8 + 2 + 1 = 1 \cdot 2^3 + 0 \cdot 2^2 + 1 \cdot 2^1 + 1 \cdot 2^0 = 1011_2$$

$$\begin{array}{r}
 4 \\
 \oplus 9 \\
 \oplus 11 \\
 \hline
 6 = 0110_2
 \end{array}$$

** Properties of nim-sum

- Commutative: $x \oplus y = y \oplus x$
- Associative: $x \oplus (y \oplus z) = (x \oplus y) \oplus z$
- $x \oplus x = 0$ for any x .
- In fact, $x \oplus y = 0$ if and only if $x = y$.

⇒ if $x \oplus y = z$ then

$y = x \oplus z$ [by adding x to both sides of \oplus]

and $x = y \oplus z$ [by adding y to both sides of \oplus]

* XOR

- $(x_1 \oplus x_2 \oplus \dots \oplus x_k)$ in binary puts a "1" in each column of the binary XOR that has an odd number of "1"s, and a "0" in each column that has an even number of "1"s.

* Theorem: Suppose that (m_1, m_2, \dots, m_k) is a game state in a game of nim.

This state is an "N" state if and only if $(m_1 \oplus m_2 \oplus \dots \oplus m_k) > 0$.

Otherwise, if the nim-sum is zero, it is a P state (and conversely).

* Examples

- Game state (m) , the nim sum is m . So it is "N" if and only if $m > 0$.

- Game state (a, b) , the nim sum is $(a \oplus b)$. It is "N" if and only if $(a \oplus b) > 0$, which occurs if and only if $a \neq b$.

- State $(1, 2, 3) \rightarrow$ it is a P state.

$$1 \oplus 2 \oplus 3 = 1_2 \oplus 10_2 \oplus 11_2 = \begin{array}{r} 1_2 \\ 10_2 \\ 11_2 \\ \hline 00_2 = 0 \end{array}$$

- State $(1, 2, 4) \rightarrow$ it is N

$$1 \oplus 2 \oplus 4 = \begin{array}{r} 1_2 \\ \oplus 10_2 \\ \oplus 100_2 \\ \hline 111_2 = 7 > 0 \end{array}$$

The theorem implies the following statements:

- ① If $m_1 \oplus \dots \oplus m_k = 0$, then any move results in a positive nim-sum - [by property of P-states]
- ② If $m_1 \oplus \dots \oplus m_k > 0$, then there is at least one move that results in a zero nim-sum. [by property/def. of N-states]

Let us check these.

- ① Suppose (m_1, \dots, m_k) is a state such that [E.g. (1, 2, 3)]
- ⊗ $m_1 \oplus \dots \oplus m_k = 0$

Suppose we eat some berries out of the first pile. The new state will be

$$(m_1', m_2, \dots, m_k)$$

The new nim-sum is

$$m_1' \oplus m_2 \oplus \dots \oplus m_k$$

Note: $(m_2 \oplus \dots \oplus m_k) = m_1$ [by adding m_1 to both sides of ⊗]

⇒ new nim-sum is

$m_1' \oplus m_1$. Since $m_1' \neq m_1$, we see that

$$m_1' \oplus m_1 = \text{new nim sum} > 0$$

The same argument works for making moves in any other pile ⇒ any move sends you to a state with a positive nim-sum.

② Let's check the other statement.

Suppose (m_1, \dots, m_k) is a state such that

$$m_1 \oplus \dots \oplus m_k > 0.$$

We need to find at least one move that makes the ~~the~~ nim-sum zero.

Let $S = m_1 \oplus \dots \oplus m_k$

Suppose we make a move in the i^{th} pile and suppose also that the new nim-sum is zero.

$$m_i \mapsto m_i'$$

$$0 = m_1 \oplus m_2 \oplus \dots \oplus m_{i-1} \oplus m_i' \oplus m_{i+1} \oplus \dots \oplus m_k$$

Adding the two equations, we see:

$$S = (m_i \oplus \dots \oplus m_k) \oplus (m_1 \oplus \dots \oplus m_i' \oplus \dots \oplus m_k)$$

$$S = m_i \oplus m_i'$$

$$\Rightarrow m_i' = m_i \oplus S$$

We just have to make sure that there is an i such that $(m_i \oplus S) < m_i$

[to be continued...]