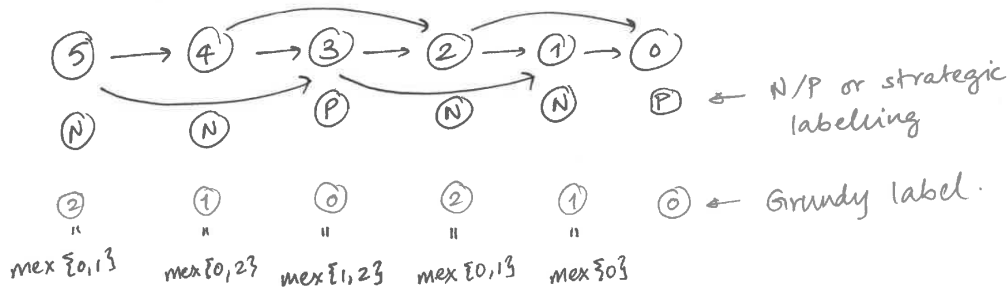


* Final exam info is on Wattle.

* Today: Grundy labelling

Example: ① subtraction game

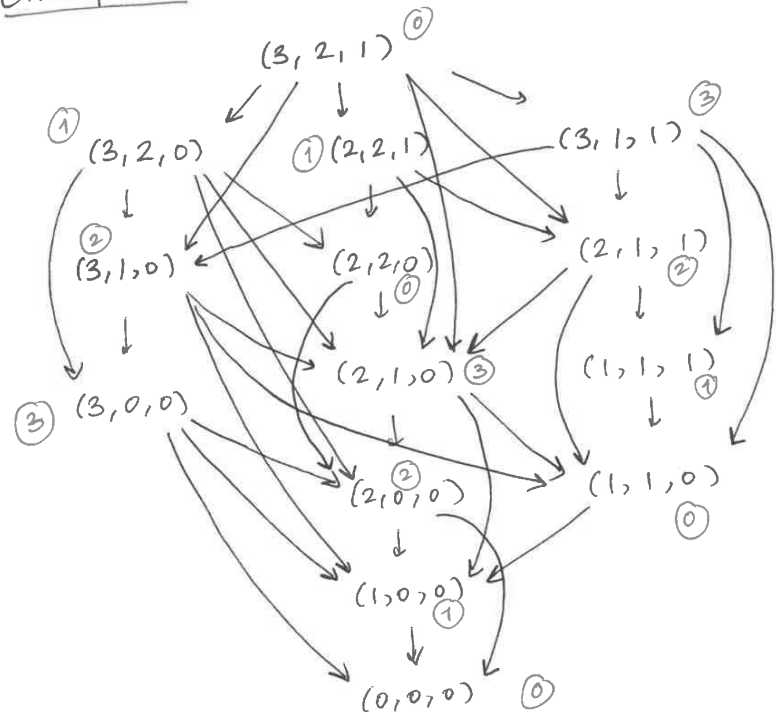
$n=5, S = \{1, 2\}$



Steps:

- ① Label terminal positions by 0
- ② Next, label any state such that ~~any~~ every place reachable from it already has a label.
- ③ Suppose that the state you're trying to label has outgoing arrows to states labelled $\{n_1, n_2, \dots, n_k\}$. This state then gets labelled by $\text{mex}\{n_1, \dots, n_k\}$
 ↑ "minimum excluded" → see examples (starting from 0)

Example 2 (Nim)



Theorem: If (m_1, \dots, m_k) is a nim game position, then its Grundy label is precisely its nim-sum $m_1 \oplus \dots \oplus m_k$.

Pf sketch: Let (m_1, \dots, m_k) be a nim position. We'll show that $s = (m_1 \oplus \dots \oplus m_k)$ is precisely the mex of all Grundy labels of states that are reachable in one step from (m_1, \dots, m_k) .
 [By induction on the graph, assume that for any state that (m_1, \dots, m_k) points to, the nim-sum equals the Grundy label.]

We need to show:

- ① No move takes us to a state with nim-sum $\neq s$.
- ② For any $s' < s$, there is a move that takes us to a state with nim-sum s' .

Let's look at ①

We're at (m_1, \dots, m_k) . If we make any move

$m_i \mapsto m'_i$, then the new nim-sum is

$$m_1 \oplus \dots \oplus m_{i-1} \oplus m'_i \oplus \dots \oplus m_k = s'$$

$$m_1 \oplus \dots \oplus m_{i-1} \oplus m_i \oplus \dots \oplus m_k = s$$

Get $m'_i \oplus m_i = s' \oplus s$.

Note: $m'_i \neq m_i \Rightarrow s' \neq s$.

Let's look at ②

Let $s' < s$, we want to find a move

$$m_i \mapsto m'_i$$

such that the new nim-sum is s' .

As before, $m'_i \oplus m_i = s' \oplus s$

$$\Rightarrow m'_i = s' \oplus s \oplus m_i$$

Need some i such that $m'_i < m_i$

$$\begin{array}{r} s_2 = 1 * \dots \boxed{1} * \dots * \\ \oplus s'_2 = 1 * \dots \boxed{0} * \dots * \\ \hline (s' \oplus s)_2 = 0 0 0 \dots \boxed{1} * \dots * \end{array}$$

← the first column from the left where $s \neq s'$ differ;
a "1" in s becomes a "0" in s' .

$$m_i = (s' \oplus s) \oplus m_i$$

\exists at least one i that has a "1" in the same column as the initial column of $(s' \oplus s)$

By the same argument as last week, there is some i for which we can make a move.

□.