

* Recap/formal def

Def: Fix an alphabet Σ . A regular expression (or regex) is a word r in the letters of Σ , together with symbols $*$, $|$, ϵ , ϕ , satisfying one of the following:

(1) $r = a$ for $a \in \Sigma$

(2) $r = \epsilon$

(3) $r = \phi$

(4) $r = r_1 r_2$ } where r_1, r_2 are valid regular expressions

(5) $r = r_1 | r_2$ ← "OR"

(6) $r = r_1^*$

* Note: Parentheses $()$ may also be used for dis-ambiguation or grouping

* As always, assume that none of the symbols $*$, $|$, ϵ , ϕ , $(,)$ are in Σ .

(2)

* Today : Matching and the language of a regexp.

Def : Given a regex r , the language of r , denoted $L(r)$, is a language on Σ ~~defined~~ defined as follows:

(1) If $r = a$ for $a \in \Sigma$, $L(r) = \{a\}$

(2) If $r = \epsilon$, $L(r) = \{\epsilon\}$

(3) If $r = \phi$, $L(\phi) = \phi$ i.e. $L(r) = \phi$

(4) If $r = r_1 r_2$, $L(r) = L(r_1) \circ L(r_2)$

(5) If $r = r_1 | r_2$, $L(r) = L(r_1) \cup L(r_2)$

(6) If $r = r_1^*$, $L(r) = L(r_1)^*$

Examples :

(1) $r = (01 | 11)^*$
 $L(r) = L((01 | 11)^*)$

Sub-question: What is:

$$L(01 | 11) ? = L(01) \cup L(11)$$

$$= \{01\} \cup \{11\} = \{01, 11\}$$

$$L((01 | 11)^*) = \{\epsilon, 01, 11, \underline{0111}, 0101, 1111, 1101, 111101, \dots\}$$

* Examples

(2) $r = (10)^* 1 \mid (01)^* 0$

$L(r) = L((10)^* 1) \cup L((01)^* 0)$

$L((10)^* 1) = L((10)^*) \circ L(1)$
 $= L(10)^* \circ \{1\}$

$= \{\epsilon, 10, 1010, 101010, \dots\} \circ \{1\}$

$= \{1, 101, 10101, 1010101, \dots\}$

$L((01)^* 0) = \{0, 010, 01010, \dots\}$

$L(r)$ = strings with at least one letter that begin & end with the same letter, and alternate letters.

* Notation: The regex $(011)^*$ has language Σ^*

Shorthand for (011) is Σ

(more generally, Σ means $(a_1 a_2 \dots a_n)$ if

$\Sigma = \{a_1, \dots, a_n\}$)

Shorthand for $(011)^*$ is Σ^*

(3) $r = (\Sigma \Sigma 0)^*$

$L(r)$ = words of length $3k$, for $k \geq 0$, such that every third letter is a 0.

