

\* Happy end of term! See you @ office hrs...  
 Check Wattle for times. We'll either be in HN4.84 or in one of the meeting rooms near the common area/kitchen of HN 4<sup>th</sup> floor.

\* Recall: Grundy( $G$ ) = 0 iff  $G$  is a P-position  
 Grundy( $G$ ) > 0 iff  $G$  is an N-position. ②

But wait, there's more!

\* Theorem: Let  $G$  &  $H$  be two games. Then...  
 the Grundy value of  $G+H$  equals the nim-sum of the Grundy values of  $G$  &  $H$

$$\text{Grundy}(G+H) = \text{Grundy}(G) \oplus \text{Grundy}(H) !$$

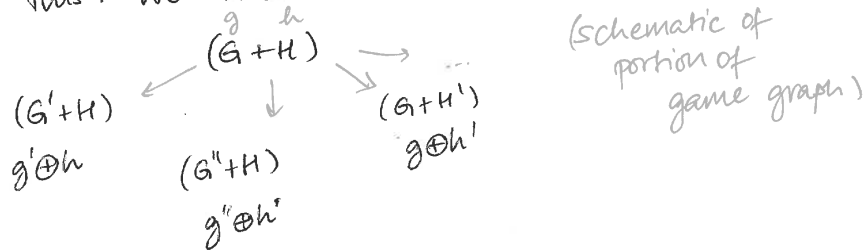
Special case: If  $G = (*m_1, \dots, *m_k)$  is a nim game, then

$$\text{Grundy}(G) = (*m_1) \oplus (*m_2) \oplus \dots \oplus (*m_k)$$

↑ game sum

Thm  $\Rightarrow$  Grundy( $G$ ) =  $(m_1 \oplus m_2 \oplus \dots \oplus m_k)$   
 (Explains labelling from Wed.)

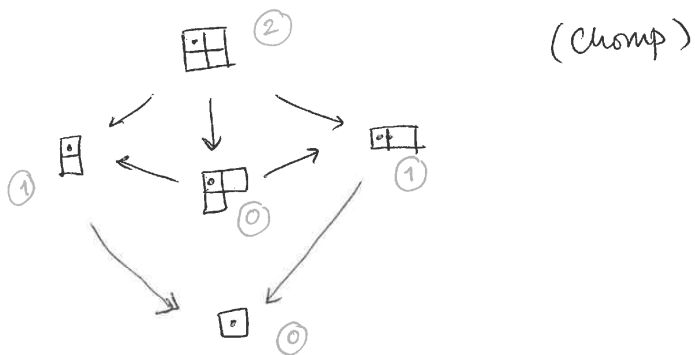
To prove this: we need to show:



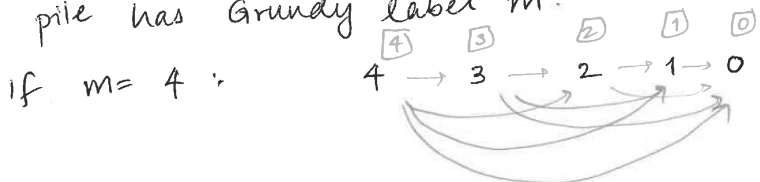
We want:  $\text{mex}\{g' \oplus h, g \oplus h'\} = g \oplus h$

\* Last time: Grundy labelling

E.g.



E.g. Single-pile nim game with  $m$  smarties in one pile has Grundy label  $m$ .



Notation: We call the single-pile nim game with  $m$  objects as  $*m$ .

So Grundy( $*m$ ) =  $m$ .

Pf sketch

Let  $g = \text{Grundy}(G)$  &  $h = \text{Grundy}(H)$

Let  $s = g \oplus h$ .

Want  $s = \text{mex} \{ g' \oplus h, g \oplus h' \mid g' = \text{Grundy}(G') \}$

(By induction, we may assume result for  $G \rightarrow G'$  valid move; for  $G'+H, G+H'$  }  $h' = \text{Grundy}(H)$  for  $H \rightarrow H'$  valid move }

That is, we need to show that:

① If we make any move in  $G+H$ , we don't get the Grundy value/nim sum to be  $s$ ,

② If  $s' < s$  then we can make a move in  $G+H$  to get to  $s'$ .

(Pf by example): Suppose  $\text{Grundy}(G) = 5$   
 $\text{Grundy}(H) = 11$

$$s = 5 \oplus 11 : \begin{array}{r} 101_2 \\ \oplus 1011_2 \\ \hline 1110_2 \end{array}$$

$$s = 14 = 1110_2$$

① Note: If we move in  $G$ , we get a Grundy value  $\neq 5$   
" " " "  $H$ , " " " "  $\neq 11$   
 $\Rightarrow$  new nim-sum after we make a move will not equal 14

②: Choose any  $s' < 14$ , e.g.  $s' = 4$ .  
We now need to find a move in  $G$  OR  $H$ , such that the new nim-sum of Grundy values equals 4.

③

$\text{Grundy}(G) = 5$

$$\begin{array}{r} 101_2 = 5 \quad g \\ \oplus 1011_2 = 11 \quad h \\ \hline 1110_2 = 14 \quad s \end{array}$$

$s' = 4 = 100_2$

Need to change  $g \rightarrow g'$  or  $h \rightarrow h'$  such that  $g' \oplus h$  (or  $g \oplus h'$ ) equals 4.

Recall:  $G$  points to states labelled 0, 1, 2, 3, 4 (maybe others, but not 5)

$H$  points to states labelled 0, 1, ..., 10 (maybe others, but not 11).

Note: Sending  $H$  to  $H'$  labelled 1 will make the new nim-sum equal to 4:

$$\begin{array}{r} 101_2 \\ 1_2 \\ \hline 100_2 = 4 \end{array}$$

Another example: Take  $s' = 9 = 1001_2$

$$\begin{array}{r} 101_2 \\ \oplus 1011_2 \\ \hline 1110_2 \quad s \end{array} \quad \begin{array}{r} 10_2 \\ 101_2 \\ \hline 1001_2 \quad s' \end{array}$$

$\hookrightarrow 1001_2 \quad s'$

Explanation (hard) exercise

Idea: Look at the first column on left where  $s$  &  $s'$  differ. Since  $s' < s$ ,  $s$  will have 1 and  $s'$  a 0. Use idea from previous nim pf to figure out how to cancel that 1

④

Theorem (Sprague-Grundy theorem):

Let  $G$  be any impartial combinatorial game.

Let  $g$  be Grundy( $G$ ).

Then  $G \sim (*g)$ .

(Motto: Every game is some form of nim!)

Pf:  $G \sim (*g)$  iff  $G + (*g)$  is a P-position.

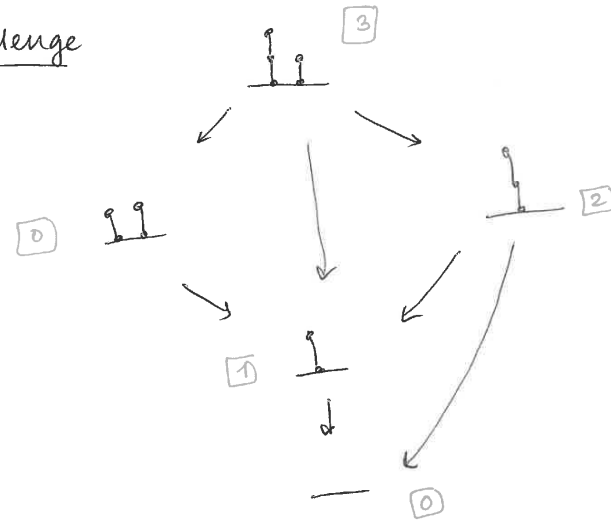
$$\begin{aligned} \text{Grundy}(G + (*g)) &= \text{Grundy}(G) \oplus \text{Grundy}(*g) \\ &= g \oplus g = 0. \end{aligned}$$

$\Leftrightarrow G + (*g)$  is a P-position. Done!

□

⑤

Eg / Challenge



⑥

~~all~~ S-G thm  $\Rightarrow$   $(3, 1, 1) \sim *3$

$\Rightarrow (3, 1, 1) + (*3)$  is a P-position.

$(3, 1, 1) + (*1)$

$(1, 1, 1) + (*1)$

$(1, 1, 0) + (*1)$

$(1, 0, 0) + (*0)$